Part III: Algebraic Types

- Lists are a special instance of an *algebraic type*.
- Algebraic types are given by a number of *constructors*, which can take parameters.
- Members of an algebraic type are accessed via *pattern matching*.
- Many functional languages provide special syntax for algebraic types and pattern matching.

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Example: Lists in Haskell (minus infix syntax, currying):

```
data List a = Nil | Cons (a, List a)
append (xs, ys) =
case xs of
Nil \Rightarrow ys
| Cons (x, xs1) \Rightarrow Cons (x, append (xs1, ys))
```

Abstract Data Types

- Algebraic types often arise naturally as implementations of abstract data types.
- An abstract data type is defined by a set of functions with their types and a set of axioms which apply to combinations of the functions.

Example: Sets of integers.

Name of abstract data type: IntSet.

Functions:

```
empty : IntSet
insert: (Int, IntSet) \rightarrow IntSet
member: (Int, IntSet) \rightarrow Boolean
```

Axioms:

```
\begin{array}{ll} \text{member } (\mathsf{x}, \ \text{empty}) &= \text{false} \\ \text{member } (\mathsf{x}, \ \text{insert } (\mathsf{x}, \ \mathsf{s})) &= \text{true} \\ \text{member } (\mathsf{x}, \ \text{insert } (\mathsf{y}, \ \mathsf{s})) &= \text{member } (\mathsf{x}, \ \mathsf{s}) \quad \  \mathbf{if} \ \mathsf{x} \ != \ \mathsf{y} \end{array}
```

Note that the behavior of the type is completely defined, even though no data representation or function implementations are given.

Reference:

J. Guttag: Abstract Data Types and the Development of Data Structures. *Communications of the ACM* vol. 20, nr. 3, pages 396–404.

IntSet's as a Data Type

An implementation of IntSets as a data type can be derived from the specification as follows:

 $\begin{array}{ll} \mbox{data IntSet} = \mbox{Empty} \mid \mbox{Insert (Int, IntSet)} \\ \mbox{member (x, Empty)} & = \mbox{false} \\ \mbox{member (x, Insert (y, s))} & = \mbox{if } (x == y) \mbox{ then true} \\ \mbox{else member (x, s)} \\ \end{array}$

Can this be generalized?

General Implementation Scheme

We can often split the set of functions defined in an abstract data type into two sets:

- A set of generators G.
- A set of accessors A.

The sets should be chosen such that each axiom is of the form

$$A(G_1, \dots, G_n) = \dots$$

If all axioms can be represented in this way, an implementation in terms of an algebraic data type suggests itself:

- Every generator function G becomes a constructor of the algebraic type.
- Every accessor function A becomes a function with all axioms starting with A as defining equations.

Encoding Algebraic Types

How can algebraic data types be represented in Funnel?

Idea: Look at Church encodings:

- A type with alternatives needs to implement the corresponding **case** expression.
- The **case** expression takes one function per alternative as parameter.
- This function represents the branch corresponding to the alternative.
- In funnel, we group the set of branch functions in a record.

Encoding Lists

```
A case expression for lists can be represented as a record:

\begin{cases}
    def Nil = ... // branch for Nil lists \\
    def Cons (x, xs) = ... // branch for Cons lists \\
    }
\end{cases}
```

Let's assume lists are objects with a match method, which takes a record representing a case expression and invokes the right branch of this record.

Then append would be coded as follows:

```
def append (xs, ys) =
    xs.match {
        def Nil = ys
        def Cons (x, xs1) = List.Cons (x, append (xs1, ys))
        }
    What does this remind you of?
```

Constructing Lists

It remains to define how lists are constructed.

As before, we will have two constructors, Nil and Cons, wrapped in a List module.

Each constructor needs to define just the match method; everything else can then be defined in terms match.

This leads to the following structure:

```
val List = {

def Nil = { def match v =

def Cons (x, xs) = { def match v =

}

How is this completed?
```

```
All other operations on lists can be written in terms of match.
For example:
      def isEmpty (xs) = xs.match {
          \mathbf{def} \operatorname{Nil} = \mathsf{true}
          \mathbf{def}\;\mathsf{Cons}\;(\mathsf{x},\,\mathsf{xs1})=\mathsf{false}
      }
      def head (xs) = xs.match {
          \mathbf{def} \operatorname{Nil} = \operatorname{error}
          def Cons (x, xs1) = x
      }
      def tail (xs) = xs.match {
          \mathbf{def} \operatorname{Nil} = \operatorname{error}
          \mathbf{def}\;\mathsf{Cons}\;(\mathsf{x},\,\mathsf{xs1})=\mathsf{xs1}
      }
Exercise: Write implementations of map, foldI and zip which use the new
representation of lists.
Exercise: Write an implementation of IntSet in Funnel.
```