Concurrency: Theory, Languages and Programming

– Pi Calculus Examples –

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$$
\begin{aligned}
B(\,i,o\,) &\stackrel{\text{def}}{=} i(x).C\langle\, x,i,o\,\rangle \\
C(\,x,i,o\,) &\stackrel{\text{def}}{=} \overline{o}\langle x\rangle.B\langle\, i,o\,\rangle \\
&\quad + \; i(y).(\; C\langle\, y,i,o\,\rangle \frown C\langle\, x,i,o\,\rangle\;)
\end{aligned}
$$

where

$$
X\langle\, i,o\,\rangle \frown Y\langle\, i,o\,\rangle \stackrel{\text{def}}{=} (\boldsymbol{\nu} m)\,(\; X\langle\, i,o\,\rangle[{}^{m}/_{o}] \mid Y\langle\, i,o\,\rangle[{}^{m}/_{i}]\;)
$$

- Follow the sequence $\xrightarrow{i1}\xrightarrow{i2}\xrightarrow{\overline{o}2}\xrightarrow{\cdots}$ to convince yourself that the buffer process is indeed a buffer (FIFO) and that it can grow unboundedly.
- Note the "type" of the stored values ...
- Note the behavior of empty cells inside a buffer "chain".

Elastic Buffers

Make the buffer elastic,
i.e., make empty cells disappear!

Several design decisions to be taken concern the question *when* an empty cell should cut itself out of a chain and die.

- if empty cell is next to a full/empty cell?
- if empty cell is left/right to a cell?
- should it be *allowed* (suicide) or *forced* (murder) to die?

One goal of this exercise is to make you think about how to argue for or against that the various design decisions above lead to equivalent solutions.

Elastic Buffers: Setup

$$
\begin{aligned}
B(\,i,l,o,r\,) &\stackrel{\text{def}}{=} i(x).C_x\langle\, i,l,o,r\,\rangle \\
&\quad + \; \ldots \\
C_x(\,i,l,o,r\,) &\stackrel{\text{def}}{=} \overline{o}\langle x\rangle.B\langle\, i,l,o,r\,\rangle \\
&\quad + \; i(y).(\; C_y\langle\, i,l,o,r\,\rangle \frown C_x\langle\, i,l,o,r\,\rangle\;) \\
&\quad + \; \ldots
\end{aligned}
$$

where

$$
\begin{aligned}
(\; X \frown Y\;)\langle\, i,l,o,r\,\rangle &\stackrel{\text{def}}{=} \\
X\langle\, i,l,o,r\,\rangle \frown Y\langle\, i,l,o,r\,\rangle &\stackrel{\text{def}}{=} \\
&\quad \ldots
\end{aligned}
$$

$$
\begin{array}{rcl}
B & \stackrel{\text{def}}{=} & (i,l,o,r). \\
 & & \quad i(x).C_x\langle\, i,l,o,r\,\rangle \\
 & & +\ \ldots \\
C_x & \stackrel{\text{def}}{=} & (i,l,o,r). \\
 & & \quad \overline{o}\langle x\rangle.B\langle\, i,l,o,r\,\rangle \\
 & & +\ i(y).\big(\ (C_y \frown C_x)\langle\, i,l,o,r\,\rangle\ \big) \\
 & & +\ \ldots
\end{array}
$$

$$
\begin{array}{rcl}
B & \stackrel{\text{def}}{=} & (i,l,o,r). \\
 & & \quad i(x).C_x\langle\, i,l,o,r\,\rangle \\
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 & & +\ i(y).\big(\ (C_y \frown C_x)\langle\, i,l,o,r\,\rangle\ \big) \\
 & & +\ \ldots
\end{array}
$$